

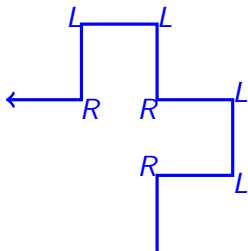
Codings of rotations are full

A. Blondin Massé S. Brlek S. Labbé L. Vuillon

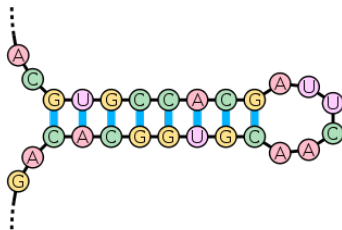
Université du Québec à Montréal

EUROCOMB 2009
September 11th, 2009

Palindromes

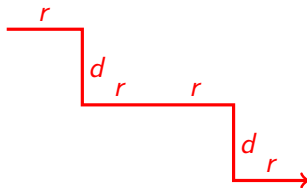


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The Fibonacci word

We define $f_{-1} = b$, $f_0 = a$ and, for $n \geq 1$,

$$f_n = f_{n-1}f_{n-2}.$$

Therefore, we have

$$f_0 = a$$

$$f_1 = ab$$

$$f_2 = aba$$

$$f_3 = abaab$$

$$f_4 = abaababa$$

$$f_5 = abaababaabaab$$

$$\vdots \quad \vdots$$

The infinite word f_∞ is called the **Fibonacci word**.

The Thue-Morse word

We define $t_0 = a$ and, for $n \geq 1$,

$$t_n = t_{n-1} \overline{t_{n-1}}.$$

so that

$$t_0 = a$$

$$t_1 = ab$$

$$t_2 = abba$$

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$$t_4 = abbabaabbaababba$$

$$t_5 = abbabaabbaababbabaababbaabbabaab$$

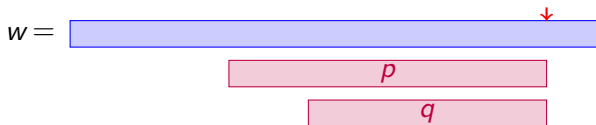
$$\vdots \quad \vdots$$

The infinite word t_∞ is called the **Thue-Morse word**.

Number of distinct palindromic factors

Theorem (Droubay, Justin and Pirillo, 2001)

Let w be a finite word. Then $|\text{Pal}(w)| \leq |w| + 1$.

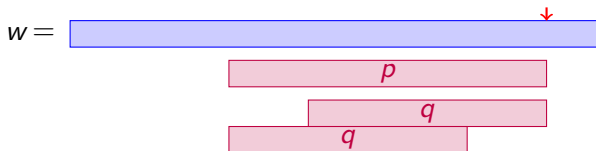


- Assume that the **first** occurrence of some palindromes p and q **ends** at the same position.

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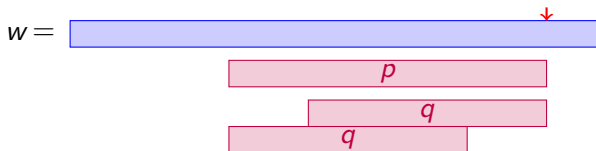


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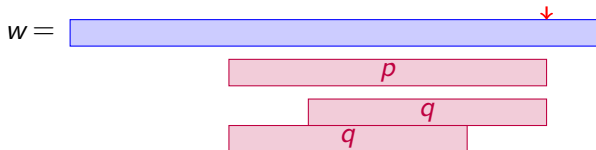


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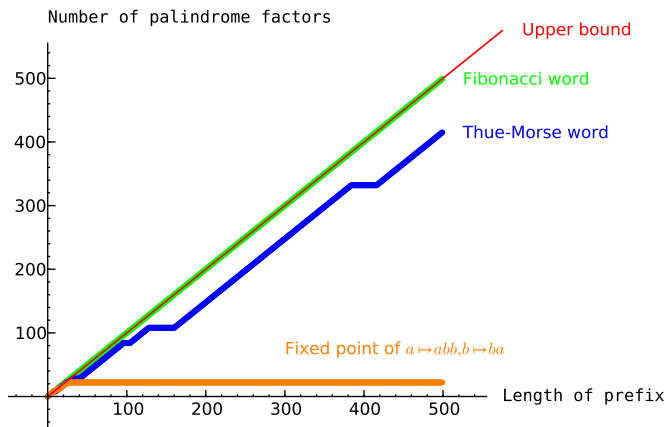


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- Then $p = q$.

Theorem (Droubay, Justin and Pirillo, 2001)

Sturmian words are *full*, i.e. they realize the *upper bound*.

Palindromic complexity



The Fibonacci word is full

$$w = a$$

Palindromes a

The Fibonacci word is full

$$w = a \ b$$

Palindromes a
 b

The Fibonacci word is full

$$w = a \ b \ a$$

Palindromes

<i>a</i>
<i>b</i>
<i>a b a</i>

The Fibonacci word is full

$$w = a \ b \ a \ a$$

Palindromes

a
 b
 $a \ b \ a$
 $a \ a$

The Fibonacci word is full

$$w = a \ b \ a \ a \ b$$

Palindromes

a
b
a b a
a a
b a a b

The Fibonacci word is full

$w = a b a a b a$

Palindromes

a
 b
 $a b a$
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The Fibonacci word is full

$w = a b a a b a b$

Palindromes

a
 b
 $a b a$
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Palindromes

a
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The Fibonacci word is full

$w = a b a a b a b a a \dots$

Palindromes

a
 b
 $a b a$
 $a a$
 $b a a b$
 $a b a a b a$
 $b a b$
 $a b a b a$
 $a a b a b a a$
 \dots

The Thue-Morse word is lacunary

$$w = a$$

Palindromes a

The Thue-Morse word is lacunary

$$w = a \mathit{b}$$

Palindromes a
 b

The Thue-Morse word is lacunary

$$w = a \text{ } b \text{ } b$$

Palindromes

a

b

b b

The Thue-Morse word is lacunary

$$w = a \ b \ b \ a$$

Palindromes

a
 b
 $b \ b$
 $a \ b \ b \ a$

The Thue-Morse word is lacunary

$$w = a \ b \ b \ a \ b$$

Palindromes

a
 b
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Palindromes

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 $a \ a$
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The Thue-Morse word is lacunary

$w = a b b a b a a b b \dots$

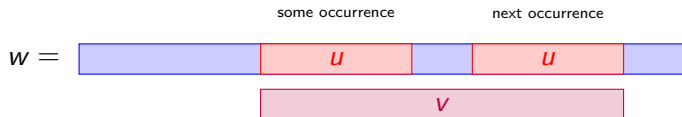
Palindromes

a
 b
 $b b$
 $a b b a$
 $b a b$
 $a b a$
 $a a$
 $b a a b$

—
...

There is **no** new palindrome at this position!

Complete return words



We say that v is a **complete return word** of u in w , if v starts at an occurrence of u and ends at the end of the next occurrence of u .

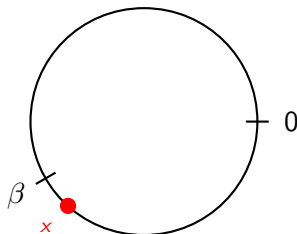
Fact

A word w is **full** if and only if every **complete return word** of a **palindrome** factor of w is a **palindrome**.

Codings of rotations (1/2)

The **coding of rotations** of parameters (x, α, β) is the word $\mathbf{C} = c_0 c_1 c_2 \dots$ such that

$$c_i = \begin{cases} 0 & \text{if } x + i\alpha \in [0, \beta) \\ 1 & \text{if } x + i\alpha \in [\beta, 1) \end{cases}$$

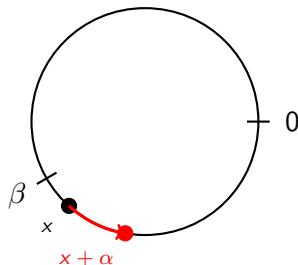


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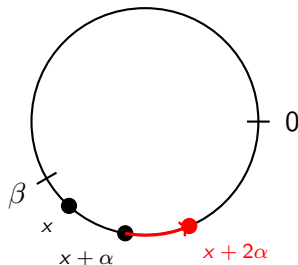


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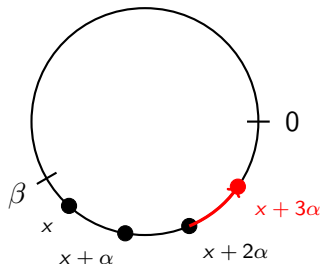


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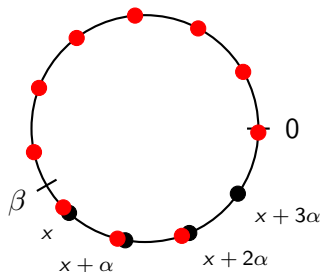


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Codings of rotations (2/2)

Many interesting problems related to codings of rotations:

- **Density** of the letters 0 and 1,
- **Complexity**, i.e. the number of factors of length n , or **palindromic** and **f -palindromic** complexity,
- Applications to **number theory** [Adamczewski, 2002],
- etc.

In particular, Rote (1994) expressed sequences of complexity $2n$ with respect to codings of rotations.

The different cases

Let \mathbf{C} be a coding of rotations of parameters (x, α, β) .

- If α is **rational**, then \mathbf{C} is **periodic**.
- If $\alpha = \beta$ is **irrational**, then \mathbf{C} is **Sturmian**

$$f(n) = n + 1.$$

- If α and β are **rationally dependent**, then \mathbf{C} is **quasi-Sturmian**.

$$f(n) = n + k, \quad \text{for some constant } k.$$

- Otherwise, \mathbf{C} is a **Rote sequence**

$$f(n) = 2n, \quad \text{for large enough } n.$$

Theorem

*Every coding of rotations is **full**.*

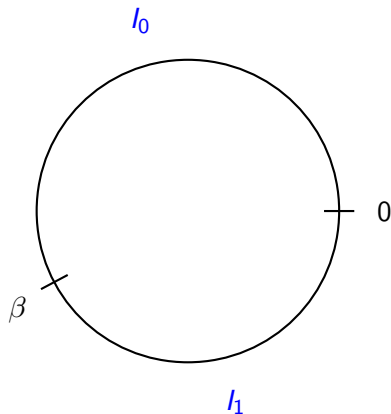
The proof is based on the following ideas:

- 1 Return words
- 2 Interval exchange transformations
- 3 Poincaré's first return function
- 4 Many results on those dynamical systems

Idea of the proof

Let $x = 0.102$, $\alpha = 0.135$ and $\beta = 0.578$. Then

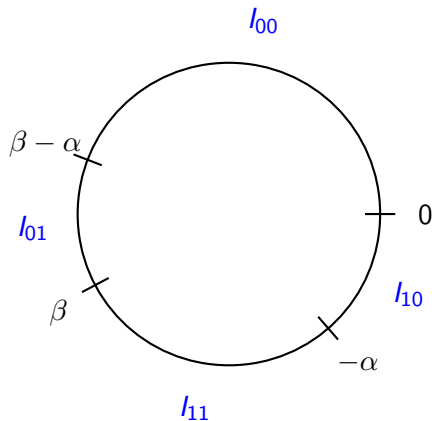
$\mathbf{C} = 0000111000011110000111000011100000111000 \dots$



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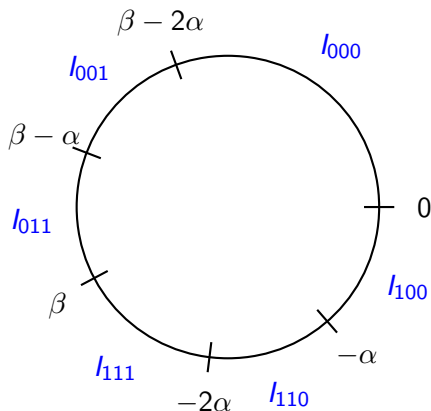
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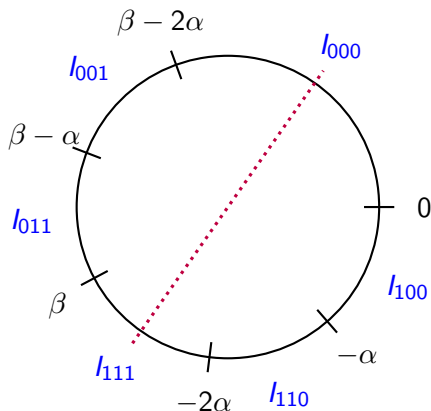
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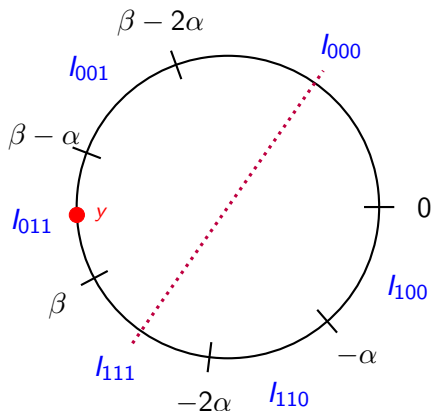
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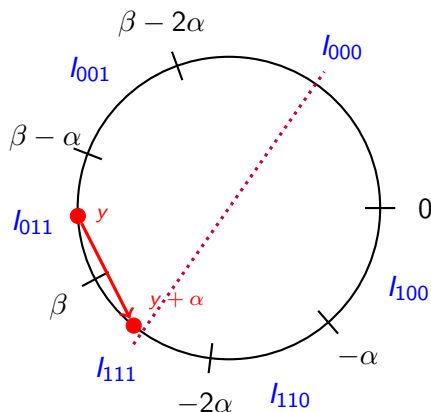
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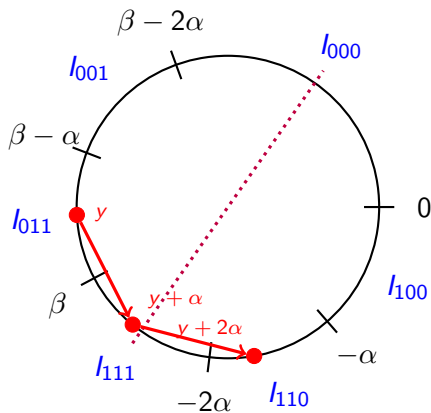
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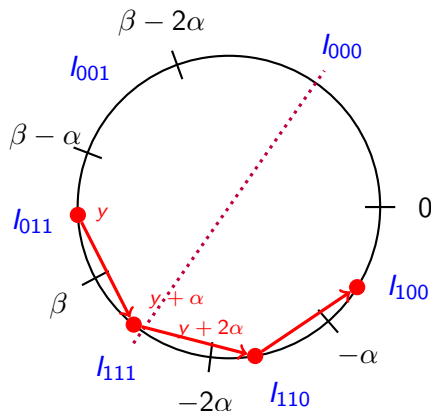
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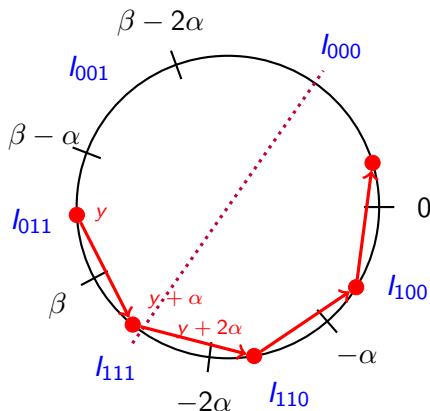
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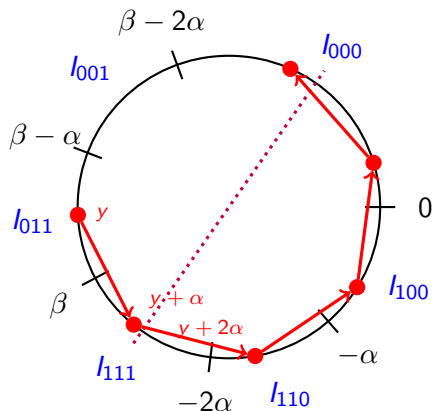
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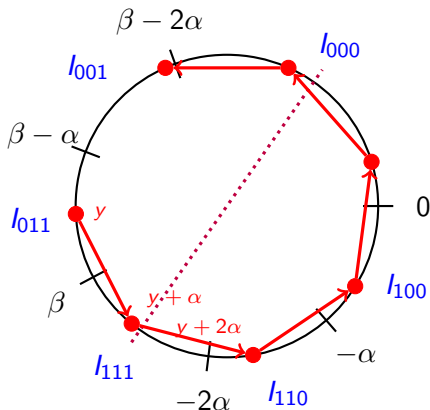
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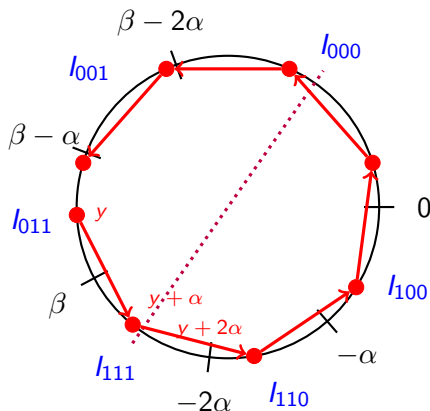
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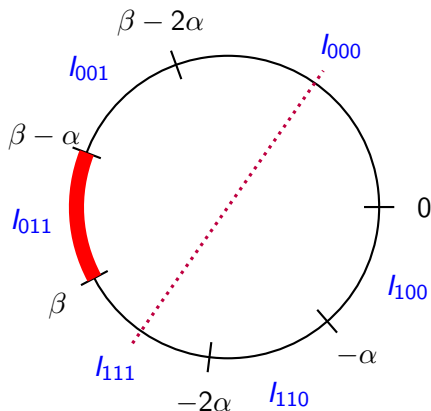
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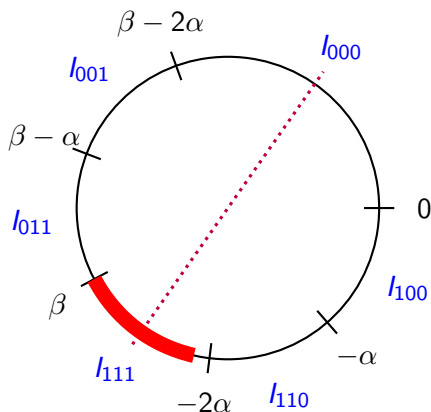
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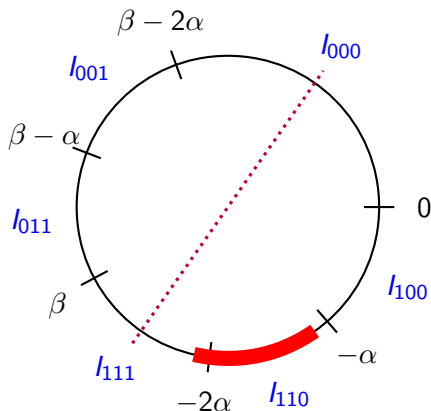
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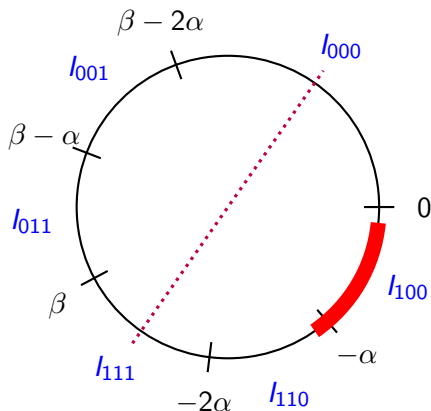
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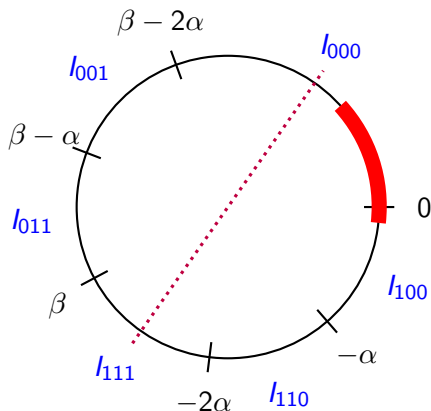
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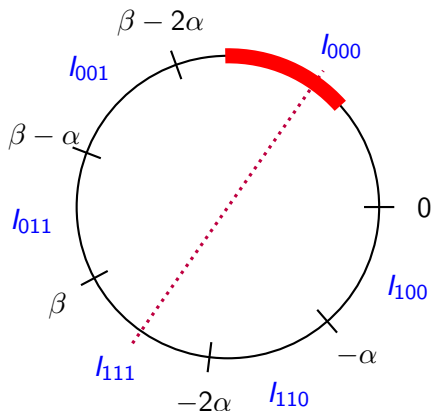
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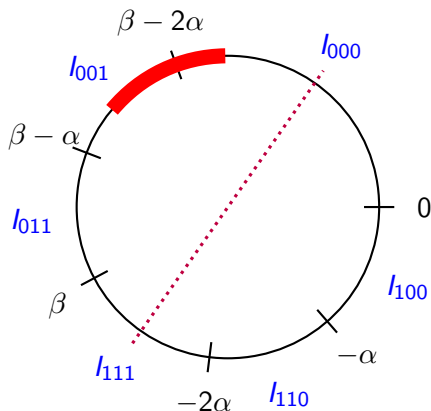
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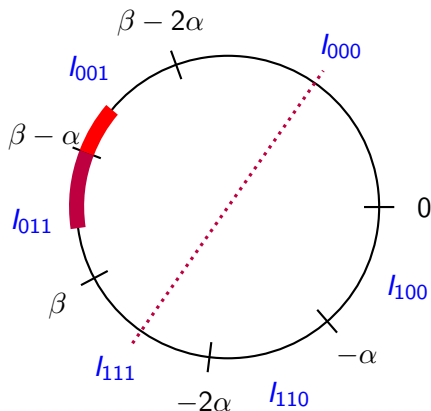
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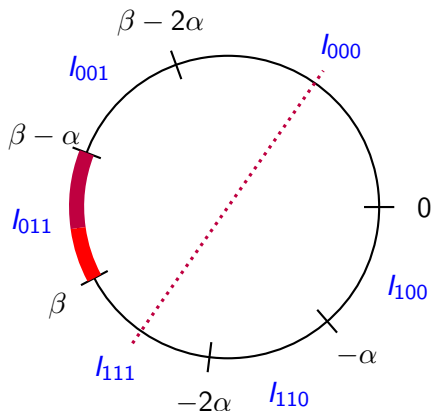
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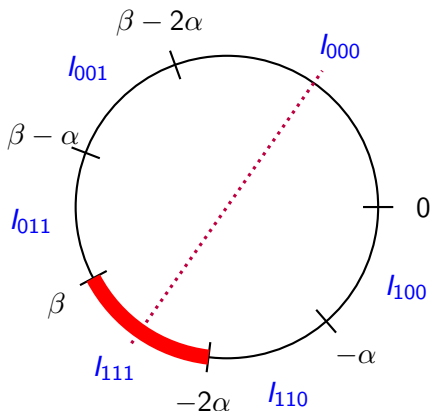
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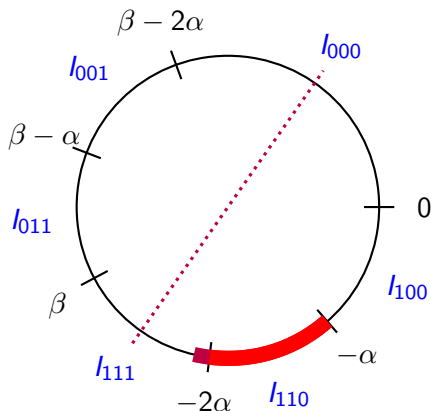
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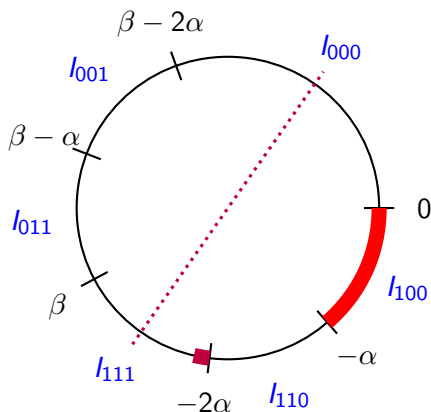
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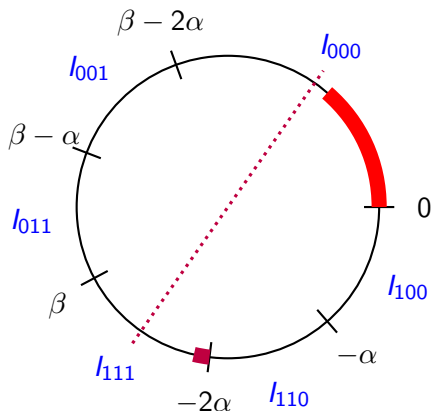
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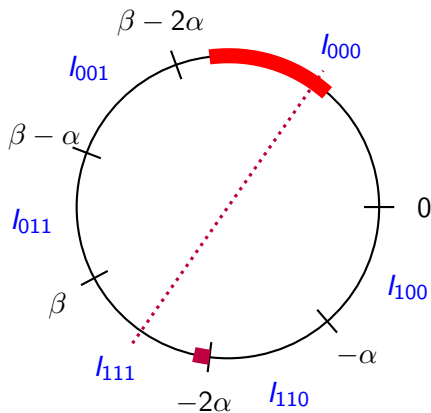
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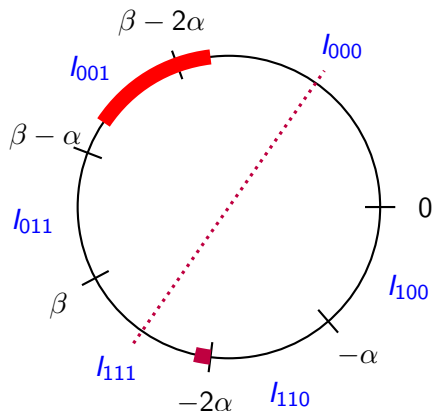
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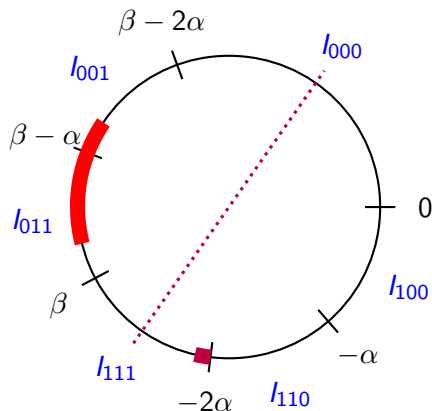
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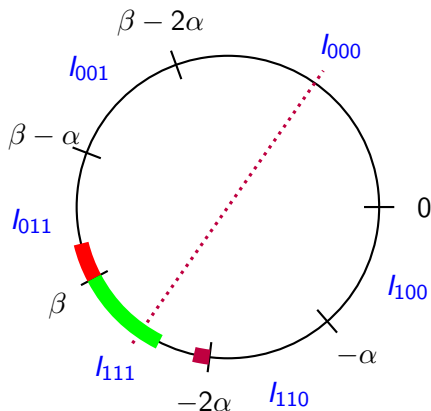
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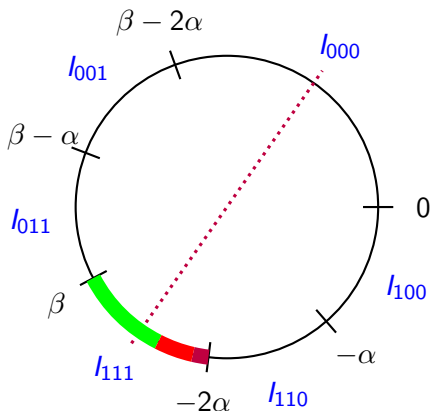
$\mathbf{C} = 0000111000011110000111000011100000111000 \dots$



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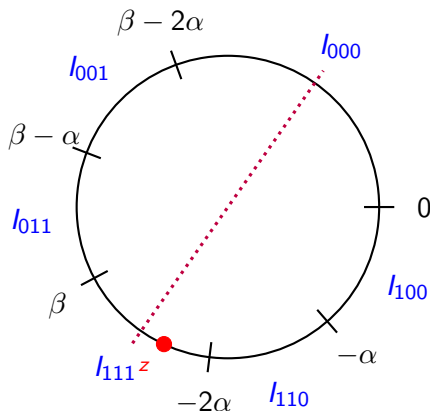
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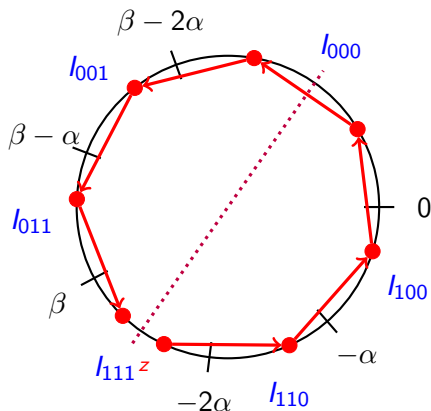
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This research was driven by computer exploration using the open-source mathematical software **Sage** [1] and its algebraic combinatorics features developed by the **Sage-Combinat** community [2], and in particular, F. Saliola, A. Bergeron and S. Labbé.

The pictures have been produced using Sage and **pgf/tikz**.



W. A. Stein et al., *Sage Mathematics Software (Version 4.1.1)*, The Sage Development Team, 2009, <http://www.sagemath.org>.



The Sage-Combinat community, Sage-Combinat: enhancing Sage as a toolbox for computer exploration in algebraic combinatorics, <http://combinat.sagemath.org>, 2009.